

FACULTY OF ENGINEERING & TECHNOLOGY

BCS-501 Operating System

Lecturer-14

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Deadlocks

- **≻System Model**
- > Deadlock Characterization
- > Resource-Allocation Graph



The Deadlock Problem

•A set of blocked processes each holding a resource and waiting to acquire a resource held by another process in the set.

Example

System has 2 tape drives.

 P_1 and P_2 each hold one tape drive and each needs another one.

Example

semaphores A and B, initialized to 1

$$P_0$$
 P_1 wait (A); wait (B); wait (A)

- •To develop a description of deadlocks, which prevent sets of concurrent processes from completing their tasks
- •To present a number of different methods for preventing or avoiding deadlocks in a computer system

System Model

System consists of resources Resource types $R_1, R_2, ..., R_m$

CPU cycles, memory space, I/O devices Each resource type R_i has W_i instances.

Each process utilizes a resource as follows:

- •request
- •use
- ·release



Deadlock Characterization

Deadlock can arise if four conditions hold simultaneously.

Mutual exclusion: only one process at a time can use a resource.

Hold and wait: a process holding at least one resource is waiting to acquire additional resources held by other processes.

No preemption: a resource can be released only voluntarily by the process holding it, after that process has completed its task.

Circular wait: there exists a set {P0, P1, ..., P0} of waiting processes such that P0 is waiting for a resource that is

held by P1, P1 is waiting for a resource that is held by P2, ..., Pn–1 is waiting for a resource that is held by

Pn, and P0 is waiting for a resource that is held by P0.

Resource-Allocation Graph

•A set of vertices V and a set of edges E.

•V is partitioned into two types:

P = {P1, P2, ..., Pn}, the set consisting of all the processes in the system.

R = {R1, R2, ..., Rm}, the set consisting of all resource types in the system.

request edge – directed edge $P1 \rightarrow Rj$

assignment edge – directed edge Rj → Pi

Resource-Allocation Graph (Cont.)

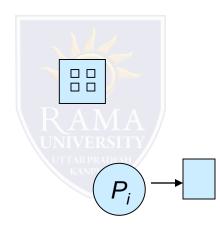
Process



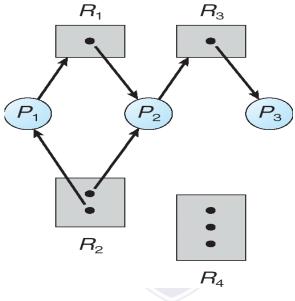
Resource Type with 4 instances

Pi requests instance of Rj

Pi is holding an instance of Rj



Example

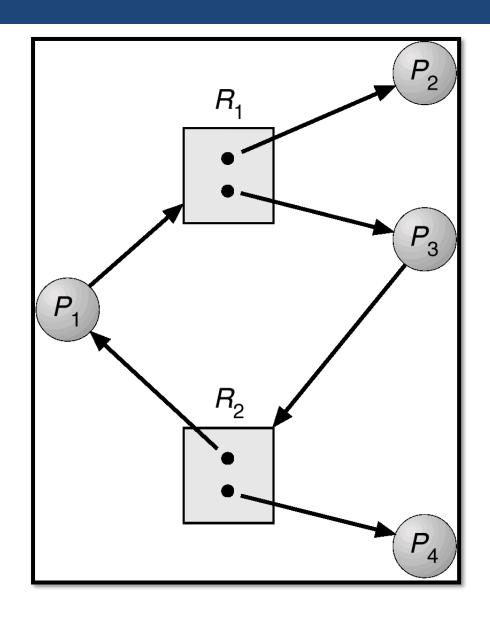


If graph contains no cycles \Rightarrow no deadlock

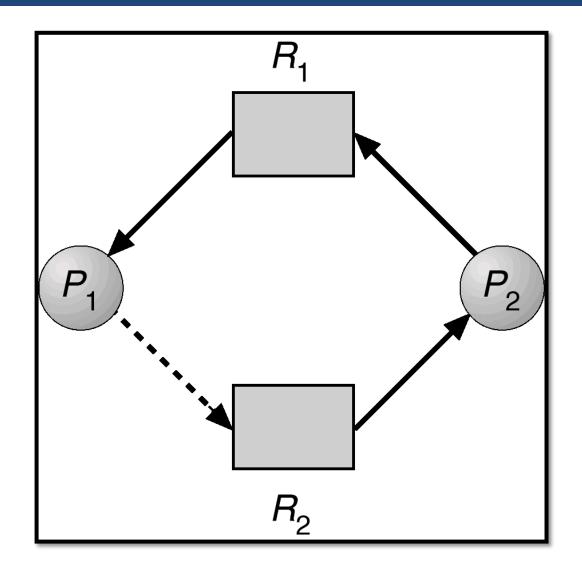
If graph contains a cycle \Rightarrow

if only one instance per resource type, then deadlock if several instances per resource type, possibility of deadlock

Resource Allocation Graph With A Cycle But No Deadlock



Unsafe State In Resource-Allocation Graph



Methods for Handling Deadlocks

- •Ensure that the system will never enter a deadlock state:
 - ➤ Deadlock prevention
 - ➤ Deadlock avoidence
- •Allow the system to enter a deadlock state and then recover

•Ignore the problem and pretend that deadlocks never occur in the system; used by most operating systems, including UNIX

MCQ

For effective operating system, when to check for deadlock?

- A. every time a resource request is made
- B. at fixed time intervals
- C. both (a) and (b)
- D. none of the mentioned

A problem encountered in multitasking when a process is perpetually denied necessary resources is called:

- A. deadlock
- B. starvation
- C. inversion
- D. aging

Which one of the following is a visual (mathematical) way to determine the deadlock occurrence?

- A. resource allocation graph
- B. starvation graph
- C. inversion graph
- D. none of the mentioned

To avoid deadlock:

- A. there must be a fixed number of resources to allocate
- B. resource allocation must be done only once
- C. all deadlocked processes must be aborted
- D. inversion technique can be used

The number of resources requested by a process:

- A. must always be less than the total number of resources available in the system
- B. must always be equal to the total number of resources available in the system
- C. must not exceed the total number of resources available in the system
- D. must exceed the total number of resources available in the system